

Graph and Digraph Glossary

[A](#) | [B](#) | [C](#) | [D](#) | [E](#) | [F](#) | [G](#) | [H](#) | [I](#) | [J](#) | [K](#) | [L](#) | [M](#) | [N](#) | [O](#) | [P-Q](#) | [R](#) | [S](#) | [T](#) | [U](#) | [V](#) | [W-Z](#)

Acyclic Graph

A [graph](#) is acyclic if it contains no [cycles](#).

Adjacency Matrix

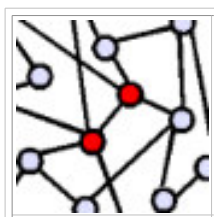
A 0-1 square matrix whose rows and columns are indexed by the vertices. A 1 in the ij -th position of the matrix means that there is an [edge](#) (or [arc](#)) from vertex i to vertex j . A 0 indicates that there is no such edge (or arc). Can be used for both graphs and digraphs.

Adjacency Structure

A representation of a graph or digraph which lists, for each vertex, all the vertices that are [adjacent](#) to the given vertex.

Adjacent

Two [vertices](#) are adjacent if they are connected by an [edge](#). We often call these two vertices *neighbors*. Two adjacent vertices:



Two edges are adjacent if they have a vertex in common.

Ancestor

In a rooted tree, a vertex on the path from the root to the vertex. Vertex v is an ancestor of vertex w if and only if w is a descendant of v .

Arc

A directed edge of a digraph. Some authors use it as a synonym for an [edge](#) of a graph. Other synonyms for arc in a digraph are arrow, directed line, directed edge, and directed link.

Arc List

A representation of a [digraph](#) using the [arcs](#) of the digraph. Can be an unordered listing of the ordered pairs, or a pair of ordered lists with the starting [vertex](#) in one list and the ending vertex in the corresponding position of the second list.

Bipartite Graph

A [graph](#) is bipartite if the [vertices](#) can be partitioned into two sets, X and Y , so that the only [edges](#) of the graph are between the vertices in X and the vertices in Y . [Trees](#) are examples of bipartite graphs. If G is bipartite, it is usually denoted by $G = (X, Y, E)$, where E is the edge set.

Binary Code

An assignment of symbols or other meanings to a set of bitstrings.

Binary Search Tree

A [binary tree](#) that has been labelled with numbers so that the right [offspring](#) and all of its descendants have labels smaller than the label of the vertex, and the left offspring and all its descendants have labels larger than that of the vertex. .

Bridge

An edge in a graph whose removal (leaving the vertices) results in a [disconnected graph](#). Also known as a *cut-edge*.

Child

In a rooted tree, a vertex v is a child of vertex w if v immediately succeeds w on the path from the root to v . Vertex v is a child of w if and only if w is the parent of v .

Chromatic Number

The chromatic number of a [graph](#) is the smallest k for which the graph is [k-colorable](#). The chromatic number of the graph G is denoted by $X(G)$. [X is the greek letter chi].

Clique

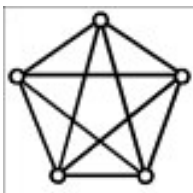
A [subgraph](#) that is a [complete](#) graph.

Closure

The closure of a graph G with n vertices, denoted by $c(G)$, is the graph obtained from G by repeatedly adding edges between non-[adjacent](#) vertices whose [degrees](#) sum to at least n , until this can no longer be done. Several results concerning the existence of [hamiltonian](#) circuits refer to the closure of a graph.

Complete

A complete graph is a simple graph in which all pairs of vertices are [adjacent](#). They are denoted by K_n , where n is the number of vertices. (The K is in honor of Kuratowski, a pioneer in graph theory.) The corresponding concept for digraphs is called a complete symmetric digraph, in which every **ordered** pair of vertices are joined by an arc. Here is the complete graph on five vertices, K_5 :



Connected Component

In a graph, a (connected) component is a maximal, [connected](#), [induced subgraph](#). Maximal means that there is no larger connected, induced subgraph containing the vertices of the component.

Condensed Graph

Given a [graph](#) G , if two [vertices](#) of G are identified and any [loops](#) or multiple [edges](#) created by this identification removed, the resulting graph is called the *condensed graph*.

Connected

A connected graph is one in which every pair of vertices are joined by a [walk](#). A graph which is not connected is called *disconnected*, and breaks up into [connected components](#).

Cycle

A closed [path](#) with at least one edge.

Decision Tree

A [binary tree](#) used to represent an algorithm for sorting by comparisons. The [leaves](#) of the tree represent the possible outcomes (orderings), while the other vertices represent test questions which have a yes or no answer.

Degree

The degree of a [vertex](#) is the the number of proper edges incident with the vertex plus twice the number of self-loops at the vertex. The degree of a graph is the maximum degree of all of its vertices.

Degree Sequence

The degree sequence of a graph is the sequence formed by arranging the vertex degrees in non-decreasing order.

Descendant

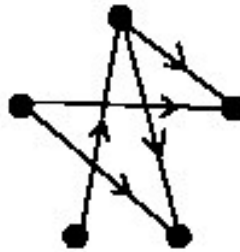
In a rooted tree, a descendant of vertex v is any vertex w whose path from the root contains v .

Diameter

The diameter of a graph is the length of the longest [walk](#) you are forced to use to get from one vertex to another in that graph. You can find the diameter of a graph by finding the [distance](#) between every pair of [vertices](#) and taking the maximum of those distances.

Digraph

A digraph is a [graph](#) in which the edges are directed and called [arcs](#). More formally, a digraph is a set of [vertices](#) together with a set of ordered pairs of the vertices, called arcs. Here is a digraph on 5 vertices:

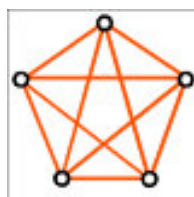


Distance

The distance between two vertices is the length of the shortest [walk](#) between them.

Edge

An edge connects two vertices in a [graph](#). We call those two vertices the endpoints of the edge. Other synonyms for edge are [arc](#), link and line. Here are the edges of a graph (in red):



Forest

A [graph](#) which contains no [cycles](#). The [connected components](#) of a forest are [trees](#).

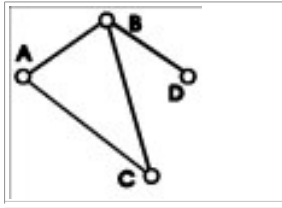
Graph

A graph is basically a collection of dots, with some pairs of dots being connected by

lines. The dots are called vertices, and the lines are called edges.

More formally, a graph is two sets. The first set is the set of vertices. The second set is the set of edges. The vertex set is just a collection of the labels for the vertices, a way to tell one vertex from another. The edge set is made up of unordered pairs of vertex labels from the vertex set.

Here is a diagram of a graph, and the sets that the graph is made from:

| | |
|---|--|
|  | $V = \{A, B, C, D\}$ --The vertex set. $E = \{(A, B), (A, C), (B, C), (B, D)\}$ --The edge set. |
| A graph diagram. | The sets that make up a graph. |

Hamiltonian

A [walk](#) or [circuit](#) in a graph is said to be *hamiltonian* if each vertex of the graph appears in it precisely once. [Paths](#) and [cycles](#) of digraphs are called hamiltonian if the same condition holds. A graph containing a hamiltonian circuit, or a digraph containing a hamiltonian cycle is referred to as a *hamiltonian graph or digraph*.

Height

The height of a [rooted tree](#) is the length of the longest [path](#) starting at the [root](#) of the tree.

Homeomorphic

Two [graphs](#) are homeomorphic if they can both be obtained from a common graph by a sequence of replacing [edges](#) by [simple chains](#). In appearance, homeomorphic graphs look like ones that have extra vertices added to or removed from edges.

Incidence Matrix

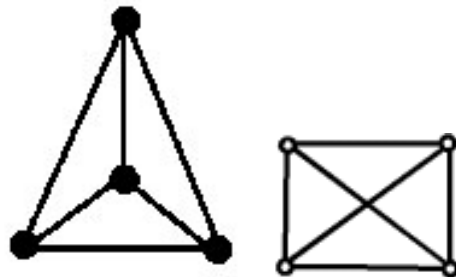
A 0-1 matrix whose rows are indexed by the vertices of a graph and whose columns are indexed by the edges. A 1 in the ij -th position of the matrix means that the vertex i is on the edge j . A 0 indicates that it is not. In some treatments, a self-loop at vertex i is indicated by a 2 in the ii -th position.

Internal Vertex

A vertex in a tree which is not a [pendant](#) vertex.

Isomorphic

Two [graphs](#) are isomorphic if they are the same graphs, drawn differently. Two graphs are isomorphic if you can label both graphs with the same [labels](#) so that every vertex has exactly the same [neighbors](#) in both graphs. Here are two isomorphic graphs:



k-Colorable

A [graph](#) is said to be k -colorable if each of its [vertices](#) can be assigned one of k colors in such a way that no two [adjacent](#) vertices are assigned the same color. The assignment is called a *coloring*.

Label

Labels are just the names we give [vertices](#) and [edges](#) so we can tell them apart. Usually, we use the integers $1, 2, \dots, n$ as the labels of a graph or digraph with n vertices. The assignment of label to vertex is arbitrary.

Leaf

A [vertex](#) of [degree](#) 1. Also known as a [pendant vertex](#).

Level

In a [rooted tree](#), the vertices at the same distance from the [root](#) are said to be at the same *level*. The root is considered to be at level 0 and the [height](#) of the tree is the maximum level.

Loop

An [edge](#) or [arc](#) from a [vertex](#) to itself is called a *loop*. Loops are not allowed in simple [graphs](#) or [digraphs](#). Also called *self-loops*.

m-ary Tree

A [rooted tree](#) in which every vertex has m or fewer [offspring](#). When $m = 2$, these are called *binary trees*. An m -ary tree is *complete* if every internal vertex has exactly m children, and all leaves have the same depth.

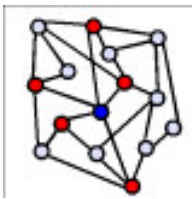
Matching

A matching in a graph is a set of edges such that every vertex of the graph is on at most one edge in the set.

Neighborhood

The neighborhood of a [vertex](#) is all the vertices that it is adjacent to (all of the

vertex's neighbors). Here we have a vertex (in blue) and the vertices in its neighborhood (in red):



Node

Another word for [vertex](#).

Offspring

In a [rooted tree](#), the vertices [adjacent](#) to a given vertex at the next higher [level](#) are called the *offspring* of the given vertex. They are sometimes called *children*. The *descendants* of a vertex are the vertices in the set of vertices which are offspring, or offspring of offspring, etc. of the given vertex..

Order

The order of a graph is the number of [vertices](#) it has.

Ordered Tree

A [rooted tree](#) in which the [children](#) of each vertex are assigned a fixed ordering.

Orientation

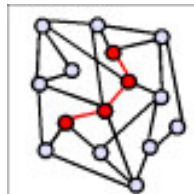
An assignment of a direction to each edge of a graph. A graph which has been given an orientation is called an *oriented graph*, and is a [digraph](#).

Parent

In a rooted tree, vertex w is the parent of vertex v if w immediately precedes v on the path from the root to v . Vertex w is the parent of v if and only if v is a child of w .

Path

A path is a [trail](#) with no repeated vertices (except possibly the first and last). Here is an example of a path:



Pendant Vertex

A [vertex](#) of [degree](#) 1. Also known as a [leaf](#).

Prefix Code

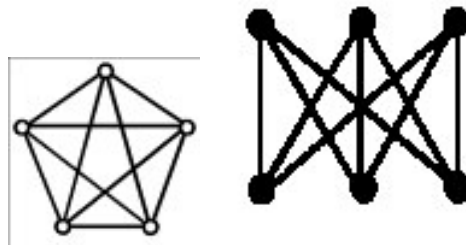
A [binary code](#) with the property that no codeword is an initial substring of any other codeword.

Perfect Matching

In a graph with $2n$ vertices, a [matching](#) with n edges is said to be perfect. Every vertex of the graph is saturated by a perfect matching. Another term for a perfect matching is a *1-factor*.

Planar

A planar graph is a graph that you can draw on a flat surface, or plane, without any of the [edges](#) crossing. Graphs that cannot be drawn on the plane without crossed edges are called non-planar graphs. Any graph that has either of the following graphs as subgraphs are non-planar:



Reduced Graph

If an [edge](#), a , is removed from a given [graph](#) G , the resulting graph, denoted G'_a is referred to as a *reduced graph*.

Regular

In a regular graph, each [vertex](#) has the same [degree](#). If this common degree is k , then we say that the graph is *k-regular*.

Rooted Tree

A [tree](#) in which one vertex has been distinguished. The distinguished vertex is called the *root* of the tree. If the tree is directed, there is a directed path from the root to each vertex of the tree.

Saturated vertex

A vertex in a graph which is on an edge of a [matching](#) is said to be *saturated*. Given a matching M , if X is a set of vertices saturated by M , then M is said to be an

X-saturating matching.

Sibling

Two vertices in a rooted tree are siblings if they have the same [parent](#).

Size

The size of a graph is the number of [edges](#) it has.

Spanning Subgraph

A [subgraph](#) of the [graph](#) G which contains all of the [vertices](#) of G .

Spanning Tree

A [spanning subgraph](#) of a [graph](#) which is also a [tree](#).

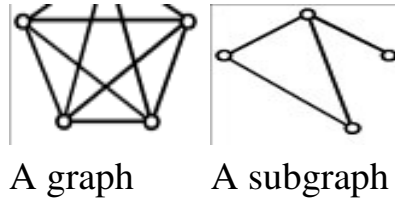
Strongly Connected

In a [digraph](#) there are many degrees of connectedness. A strongly connected digraph is one in which any vertex can be reached from any other vertex by a directed [walk](#).

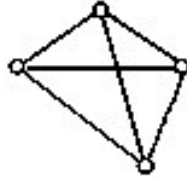
Subgraph

A subgraph of a graph is some smaller portion of that graph. Here is an example of a subgraph:





An *induced (generated) subgraph* is a subset of the vertices of the graph together with **all** the edges of the graph between the vertices of this subset. The induced subgraph of the above example is:



Topological Order

A topological ordering of a [digraph](#) is a [labelling](#) of the vertices with consecutive integers so that every [arc](#) is directed from a smaller label to a larger label.

Tournament

A tournament is a [digraph](#) in which there is exactly one [arc](#) between any two vertices. A tournament is said to be *transitive* if whenever (a,b) and (b,c) are arcs of the tournament, then (a,c) is also an arc.

Trail

In a graph, a trail is a [walk](#) with no repeated edges.

Tree

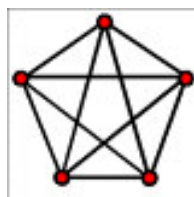
A [connected graph](#) containing no [cycles](#).

Underlying Graph

The graph that results from removing the directions on all the arcs of a digraph or partially directed graph.

Vertex

A vertex is a 'dot' in a [graph](#). The plural of vertex is 'vertices', as in, 'this graph has five vertices'. Other synonyms for vertex are [node](#), point or state. Here are the vertices of a graph (in red):



Walk

In a graph, a walk from vertex v_0 to vertex v_n is an alternating sequence

$$W = \langle v_0, e_1, v_1, e_2, \dots, v_{n-1}, e_n, v_n \rangle$$

of vertices and edges, such that the endpoints of edge e_i are v_{i-1} and v_i , for $i = 1, \dots, n$. The *length* of a walk is the number of edges in it. A walk is *closed* if $v_0 = v_n$ and *open* otherwise.
